

## Computing model and dynamic graphs

**Exercise 1** Give the families of dynamic graphs that “naturally” correspond to:

1. synchronous systems over a complete network and at most  $f$  faulty senders;
2. synchronous systems over a complete network and at most  $f$  crashes;
3. asynchronous systems over a complete network and at most  $f$  crashes;
4. asynchronous systems over a complete network and at most  $f$  initial crashes. (Do not forget links are reliable.)

To what of these families of dynamic graphs does the Santoro and Widmayer’s theorem apply?

**Exercise 2** Show that the product of  $|V| - 1$  strongly connected digraphs with the same set of nodes  $V$  and a self-loop at each node is the complete digraph.

**Exercise 3** Let  $G_1 = (V, E_1)$  and  $G_2 = (V, E_2)$  be two digraphs with the same set of nodes  $V$ , and let  $G = G_1 \circ G_2$ . Assume that each node in  $V$  has at least  $n - f_1$  and  $n - f_2$  incoming neighbors in  $G_1$  and  $G_2$ , respectively.

1. Show that any node that is not a central root of  $G$  has at most  $f_2$  outgoing neighbors in  $G_1$ .
2. Use the above bound to prove that  $G$  has at least  $n - f_1 \left(1 + \frac{f_2}{n - f_2}\right)$  central roots.
3. Suppose that a dynamic graph  $\mathbb{G}$  is such that

$$\forall i \in V, \forall t \in \mathbb{N}, |\text{In}_i(t)| > |V|/2.$$

With what delay is the dynamic graph  $\mathbb{G}$  centered?

**Exercise 4** Show that if a dynamic graph  $\mathbb{G}$  on the set of nodes  $V$  is continuously rooted, then  $\mathbb{G}$  is non-split with delay  $|V| - 1$ . Compare this result with Exercise 2.

**Exercise 5** For each positive integer  $n$ , let  $\lambda(n)$  denote the integer defined by

$$2^{\lambda(n)-1} < n \leq 2^{\lambda(n)}.$$

Let  $V$  be a finite and non-empty set of cardinality  $n$ , and let  $G_1 = (V, E_1), \dots, G_{\lambda(n)} = (V, E_{\lambda(n)})$  be  $\lambda(n)$  digraphs with the same set of nodes  $V$  and a self-loop at each node. For every index  $t \in \{1, \dots, \lambda(n)\}$ , we let

$$G_t^* = G_{\lambda(n)-t+1} \circ \dots \circ G_{\lambda(n)}.$$

1. What are the digraphs  $G_1^*$  and  $G_{\lambda(n)}^*$ ?
2. Prove that if all the digraphs  $G_1, \dots, G_{\lambda(n)}$  are non-split, then every subset of  $2^t$  nodes share a common incoming neighbor in  $G_t^*$ , for all indices  $t \in \{1, \dots, \lambda(n)\}$ .
3. What property can you then derive for dynamic graphs that are continuously non-split? For dynamic graphs that are continuously rooted (use Exercise 3)?

**Exercise 6** Let  $\mathbb{G}$  be any dynamic graph., and let  $G_1 = (V, E_1), \dots, G_{\lambda(n)} = (V, E_{\lambda(n)})$  be  $\lambda(n)$  digraphs with the same sets of nodes  $V$  and a self-loop at each node. For every index  $t \in \{1, \dots, \lambda(n)\}$ , we let

$$G_t^* = G_{\lambda(n)-t+1} \circ \dots \circ G_{\lambda(n)}.$$

1. Show that the digraphs in the integral dynamic graph  $\overline{\mathbb{G}}$  eventually stabilize to  $\overline{\mathbb{G}}$ 's limit superior, i.e.,

$$\exists s, \forall t \geq s, \overline{\mathbb{G}}(t) = \overline{\mathbb{G}}(\infty).$$

2. Prove that  $\overline{\mathbb{G}}(\infty)$  is the transitive closure of  $\mathbb{G}(\infty)$ .
3. Give a characterization of the kernel of  $\mathbb{G}$  in terms of the roots of  $\mathbb{G}(\infty)$ .

**Exercise 7** Use the positive result in the FLP paper to obtain a family of dynamic graphs that can be implemented in any asynchronous system over a complete network and a minority of initial crashes.